Mixing time of random walks on networks

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Outline

- What is a random walk?
- How do we quantify a random walk?
- Fastest mixing problem
- Computational results
- Applications

mixing time \longrightarrow convergence speed

fast mixing —> fast convergence of dynamical process on network

Random walk

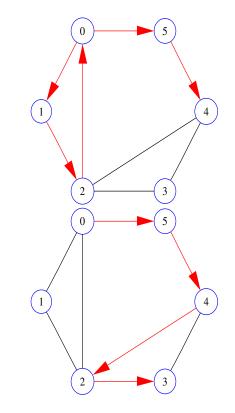
 $\blacktriangleright \Gamma(v)$: neighbors of v in G

 $x_0 \xrightarrow{\mathsf{random}} x_1 \in \Gamma(x_0) \xrightarrow{\mathsf{random}} x_2 \in \Gamma(x_1) \xrightarrow{\mathsf{random}} \cdots \xrightarrow{\mathsf{random}} x_k \in \Gamma(x_{k-1})$

Example:

 $\{v_0v_1v_2v_0v_5v_4\}$ is one instance of a random walk

 $\{v_0v_5v_4v_2v_3\}$ is another instance of a random walk



Random walk and Markov chain

 \blacktriangleright Transition probability matrix P

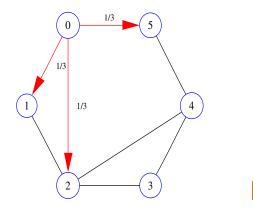
$$P_{ij} = \begin{cases} \Pr[x_{t+1} = j | x_t = i], & \text{if } ij \in E \\ 0, & \text{otherwise} \end{cases}$$

► We can choose each node uniformly

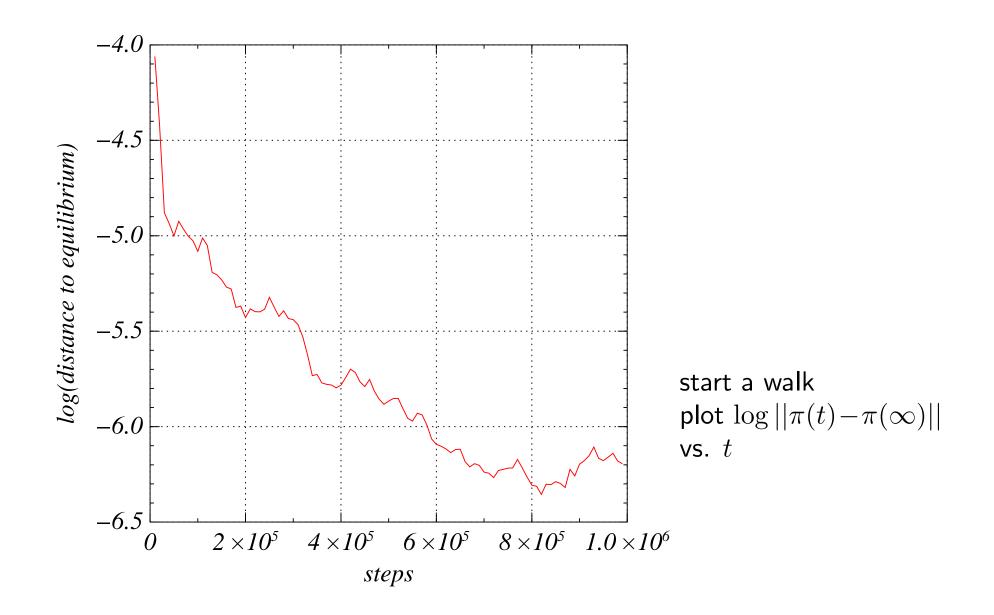
$$P_{ij} = \left\{ \begin{array}{ll} 1/d(i), & \text{if } ij \in E \\ 0, & \text{otherwise} \end{array} \right.$$

where d(i) is the degree of i

- ▶ distribution at t: $\pi(t) = \pi(0)P^t$ ■
- equilibrium distribution: $\pi = \pi P$

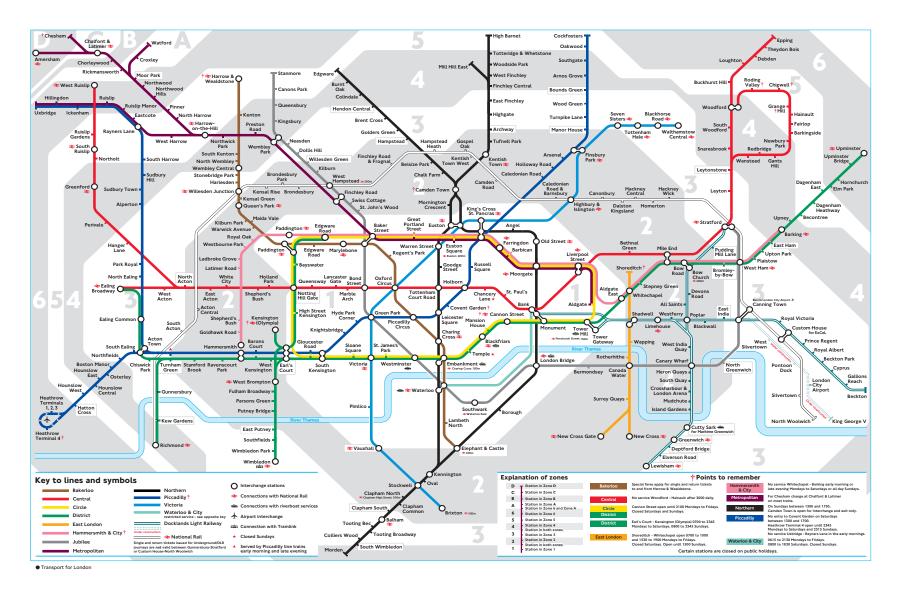


The approach to equilibrium

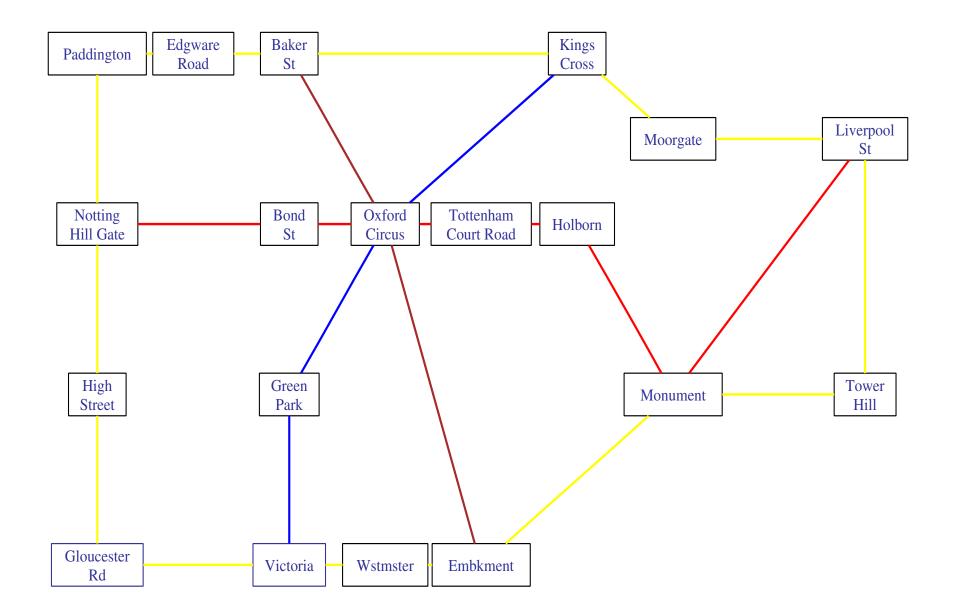


Statistical properties of random walks

Example: London underground graph



Example: London underground graph



Example: London underground graph - hitting time



Hitting time [LL00]:

$$H_{ij} = 2m \sum_{k=2}^{n} \frac{1}{1 - \lambda_k} \left[\frac{v_{ki}^2}{d(i)} - \frac{v_{ki}v_{kj}}{\sqrt{d(i)d(j)}} \right]$$

 $\{\lambda, v\}$ is the eigensystem of $N = D^{1/2}AD^{1/2}$, A is the adjacency matrix, D = diag(1/d(i)), m is the number of edges in the graph

the mean number of steps between King's Cross and Victoria is

$$H_{kc-v} = H_{81} = 38.5$$

London underground graph - Hitting time matrix

0.0	12.7	30.1	24.9	36.1	58.8	47.5	49.7	31.2	30.0	46.5	45.0	25.4	20.9	25.9	13.9	50.8	31.3	43.4
28.1	0.0	21.3	20.1	32.6	55.4	44.4	47.4	29.6	29.8	49.6	51.4	35.0	39.8	18.5	11.8	47.7	35.1	40.8
40.6	16.4	0.0	11.0	27.0	50.5	40.0	44.7	28.7	29.9	51.3	54.5	39.7	48.3	26.6	11.5	43.9	37.3	38.8
51.1	30.9	26.7	0.0	19.5	43.5	33.7	40.1	25.8	28.1	50.9	55.7	42.3	54.9	32.7	9.3	38.1	37.5	34.7
58.4	39.4	38.7	15.5	0.0	26.4	18.9	32.5	25.4	30.3	54.0	59.6	47.1	61.0	39.0	13.3	26.5	41.9	30.9
60.3	41.5	41.4	18.8	5.7	0.0	10.5	27.8	24.4	30.8	54.8	60.7	48.5	62.6	40.8	14.8	30.8	43.3	33.8
60.2	41.7	42.2	20.1	9.4	21.6	0.0	21.1	21.5	29.2	53.5	59.7	47.9	62.2	40.6	14.3	33.1	42.7	34.7
58.8	41.1	43.3	23.0	19.4	35.4	17.6	0.0	11.8	23.5	48.8	55.9	45.0	60.1	39.1	11.8	38.9	40.1	36.4
55.5	38.5	42.5	23.9	27.4	47.2	33.2	26.9	0.0	15.9	42.0	50.1	40.1	56.0	35.5	7.3	42.7	35.4	36.0
53.3	37.7	42.7	25.2	31.4	52.6	39.9	37.7	14.9	0.0	29.8	41.6	35.3	52.5	35.0	7.0	45.3	32.8	37.2
50.1	37.8	44.4	28.3	35.4	56.9	44.5	43.3	21.4	10.1	0.0	21.8	25.5	46.0	36.9	10.7	49.2	29.8	41.0
45.0	36.0	44.0	29.4	37.4	59.1	47.1	46.8	25.8	18.2	18.2	0.0	13.8	37.6	36.8	12.5	51.1	24.8	42.8
37.8	32.1	41.6	28.6	37.3	59.4	47.7	48.3	28.3	24.4	34.3	26.2	0.0	27.1	34.8	12.2	51.0	17.8	42.6
19.9	23.4	36.8	27.7	37.7	60.1	48.6	50.0	30.7	28.2	41.4	36.6	13.7	0.0	31.4	14.1	51.9	25.5	44.0
40.7	17.9	30.8	21.3	31.5	54.0	42.7	44.7	26.0	26.4	48.0	51.6	37.1	47.1	0.0	6.9	45.3	33.7	37.1
51.3	33.7	38.4	20.4	28.4	50.7	39.0	40.0	20.3	21.0	44.5	49.8	37.1	52.4	29.5	0.0	41.0	30.2	31.5
58.0	39.5	40.6	19.2	11.5	36.5	27.6	37.0	25.7	29.2	52.8	58.3	45.8	60.1	37.8	10.9	0.0	40.0	16.5
45.6	33.9	41.0	25.5	33.9	56.0	44.3	45.2	25.3	23.7	40.4	39.0	19.6	40.8	33.1	7.1	47.0	0.0	38.1
55.7	37.6	40.5	20.8	21.0	44.6	34.3	39.5	24.0	26.1	49.7	55.1	42.5	57.3	34.7	б.4	21.5	36.1	0.0

London underground graph - commute time

11 10

7

17 15 18 16

14

6

5



commute time [LL00]

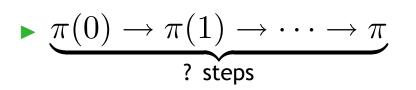
$$k_{ij} = H_{ij} + H_{ji} \quad \blacksquare$$

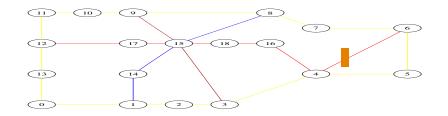
So we know that the commute time from King's Cross to Victoria then back to King's Cross is

$$k_{k-v} = k_{81} = H_{81} + H_{18} = 38.5 + 29.6 = 68.1$$

- Symmetry
 - $H_{ij} \neq H_{ji}$ unless *i* and *j* are vertex-transitive.
 - $k_{ij} = k_{ji}$ for all i, j

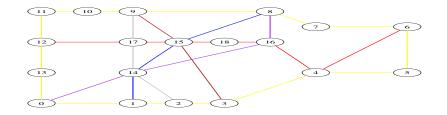
London underground graph - mixing time





- mixing rate = $\log(1/\mu(P))$ where $\mu(P) = \max_{i=2,...,n} |\lambda_i(P)| = \max\{\lambda_2(P), -\lambda_n(P)\}$
- ▶ mixing time: $\tau = 1/(\text{mixing rate}) = 1/\log(1/\mu)$ ▮
- *mixing rate* = 0.101568
 mixing time = 9.845655 Ⅰ
- build new lines

mixing rate = 0.168313mixing time = 5.941294

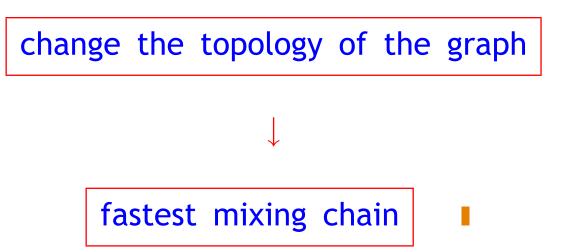


Definition of random walk properties

hitting time (access time)	H_{ij} is the expected number of steps in a random walk starting from node <i>i</i> and be- fore node <i>j</i> .	$2m\sum_{k=2}^{n}\frac{1}{1-\lambda_{k}}\left[\frac{v_{ki}^{2}}{d(i)}-\frac{v_{ki}v_{kj}}{\sqrt{d(i)d(j)}}\right]$
commute time	k_{ij} is the expected number of steps in a random walk starting at <i>i</i> , the first time return to <i>i</i> via <i>j</i> .	$H_{ij} + H_{ji}$
mixing rate	measure of how fast the random walk con- verges to its station- ary distribution.	$\rho = -\log(\mu(P))$
mixing time	the time scale (in steps) for reaching the stationary distri- bution.	$\tau = -\frac{1}{\log(\mu)}$

Fastest mixing graph problem

fastest mixing (minimum mixing time) by changing topology



Optimization description

$$\begin{array}{ll} \min_{G} & \mu(P(G)) \\ \text{s.t.} & P_{\bullet} \geqslant 0 \\ P1 = 1 \\ P_{ij} = \left\{ \begin{array}{ll} 1/d(i), & \text{if } ij \in E \\ 0, & \text{otherwise} \end{array} \right. \end{array}$$

Computational results - small regular graphs

▶ the min/max mixing time for 10 nodes regular graphs

n	deg	num	maxtime	graph	mintime	graph	avertime
10	3	17	15.5896		2.4663		6.8216
10	4	58	7.7220		1.7195		3.4093
10	5	59	7.4542		1.2427		2.2145
10	6	21	2.4663		0.9102		1.5722
10	7	5	1.1802		1.0168		1.1475

The fastest mixing problem

fastest mixing (minimum mixing time) by adjusting weights

fix the topology, changing the weights

 \downarrow (reversible chain)

fastest mixing chain

Optimization description

$$\begin{array}{ll} \min_{P} & \mu(P) \\ \text{s.t.} & P \cdot \geqslant 0 \\ & P 1 = 1 \\ & \Pi P = P^{T} \Pi \\ & P_{ij} = 0, \quad i, j \notin E \end{array}$$

Fastest mixing problem - combined problem

Fastest mixing chain both on topology and weights

all topologies

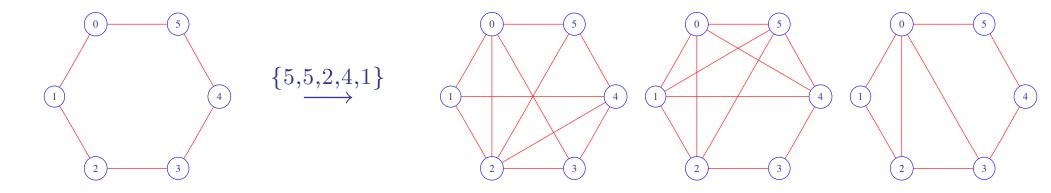


Optimization description

$$\begin{array}{ll} \min_{G} \min_{P(G)} & \mu(P) \\ \text{s.t.} & P \bullet \geqslant 0 \\ & P1 = 1 \\ & \Pi P = P^T \Pi \\ & P_{ij} = 0, \quad i, j \not\in E \end{array}$$

Small-world model $SW\{n, p\}$

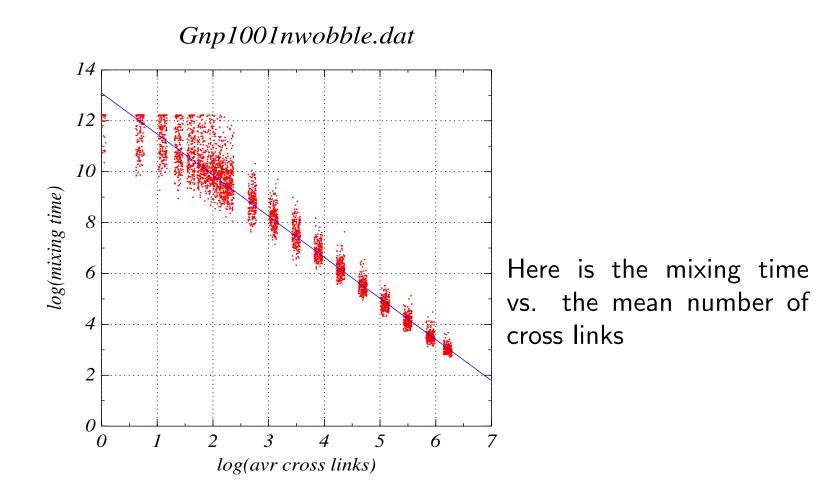
► $SW\{n, p\}$ model: n-node cycle $\Rightarrow \{m_0, m_1, \dots, m_k\} \in \text{Poisson distribution with}$ mean $\bar{m} = \binom{n}{2}p \Rightarrow SW(n, m_0), \dots, SW(n, m_k)$



 $SW\{6, 1/5\}, \bar{m} = 15 \times 0.2 = 3 \dots$

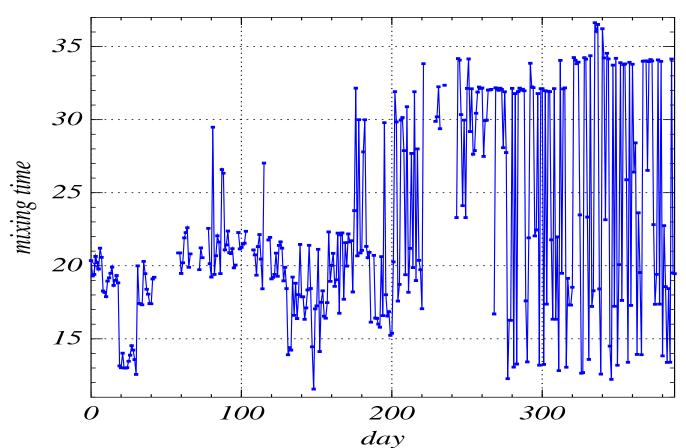
 \blacktriangleright What is the relation between \bar{m} and mixing time?

Small-world $SW\{n, p\}$ - mixing time vs. links



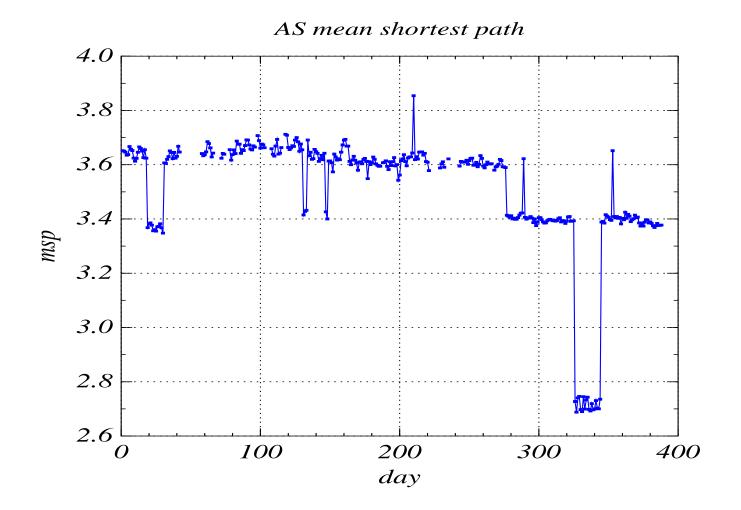
Mixing time of the internet AS graph

Internet AS graph [CAIDA] has about 10,000 nodes and changes topology constantly

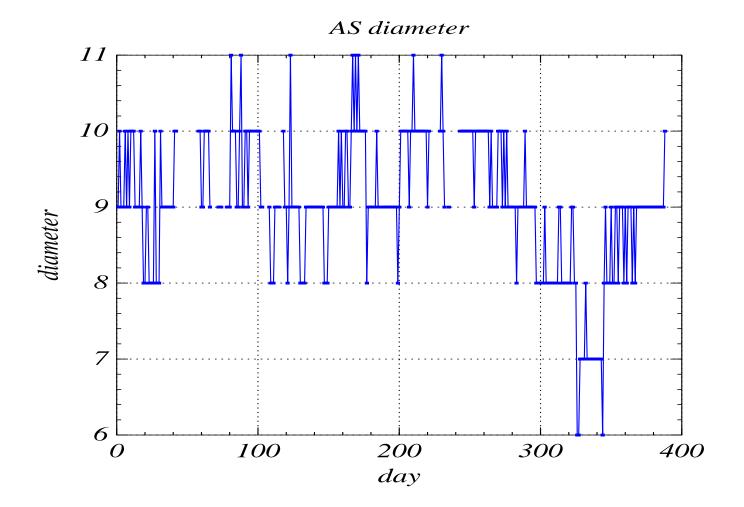


AS mixing time

Mean shortest path of the AS graph



Diameter of the internet AS graph



Applications

Information diffusion

--- How fast does information/virus/spam spread on the internet?

- Search engine (Google)
 - --- How fast can google rank the pages?
- Sampling problems
 - --- e.g. Markov chain Monte Carlo, distributed averaging

Diffusion

- **•** Eigenvalues λ_i and eigenfunctions ϕ_i
- heat kernel

$$H_t(x,y) = \sum_i \exp(-\lambda_i t) \phi_i(x) \phi_i(y)$$

satisfies

$$\frac{d}{dt}f = -\mathcal{L}_S f$$

- $\blacktriangleright H_t = \exp(-t\mathcal{L}_S)$
- this solves the diffusion problem

References

CAIDA data from sk-aslinks.caida.org/data/

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